

# Faust Term Rewriting Extension

*Albert Gräf*

*Dept. of Music Informatics*

JOHANNES  
**GUTENBERG**  
UNIVERSITÄT  
MAINZ

```
// tone.dsp

vol      = nentry("vol", 0.3, 0, 10, 0.01);
pan      = nentry("pan", 0.5, 0, 1, 0.01);
freq     = nentry("pitch", 440, 20, 20000, 0.01);

// simple sine tone generator

process = osci(freq)*vol : panner(pan);
```

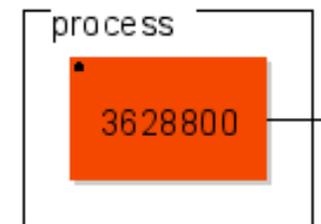
## Signal Processing with Faust

- *Functional signal processing* language, processing of *synchronous streams* of samples.
- Formal semantics means that Faust can be used as a *specification language*.
- Specifications are *executable*, sophisticated optimizations, generates competitive C++ code.
- Works with *different platforms and environments*, just recompile.

```
fact(0) = 1;  
fact(n) = n*fact(n-1);  
process = fact(10);
```

## Term Rewriting Extension

- Faust signal processors are essentially **terms** in the **block diagram algebra (BDA)**
- **Term rewriting** provides us with a means to manipulate BDA terms in an **algebraic fashion at compile time**



## Brief Digression: Term Rewriting

$\text{top}(\text{push}(s,x)) \rightarrow x$   
 $\text{pop}(\text{push}(s,x)) \rightarrow s$

term rewriting  
system

terms as “data”

reduce

$\text{top}(\text{pop}(\text{push}(\text{empty}, 1))) \rightarrow \text{top}(\text{empty})$

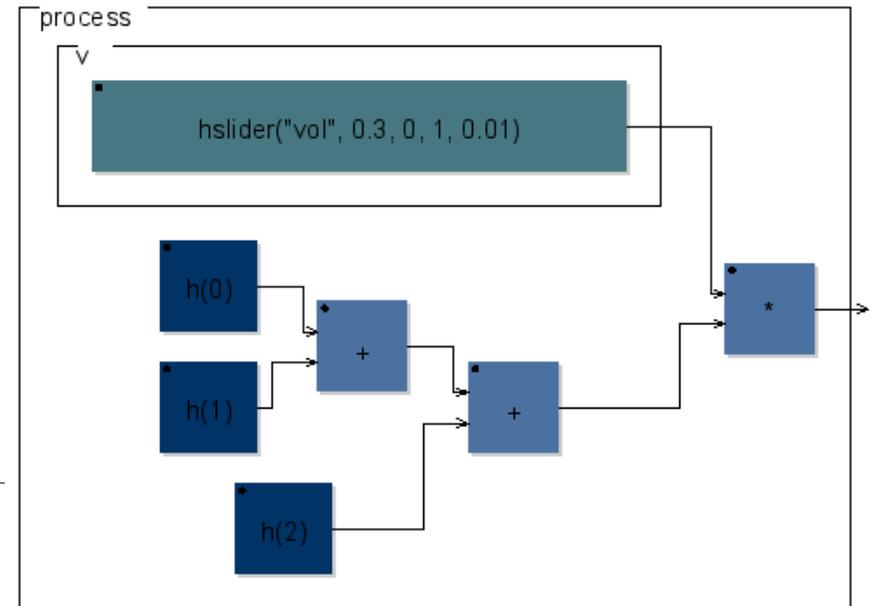
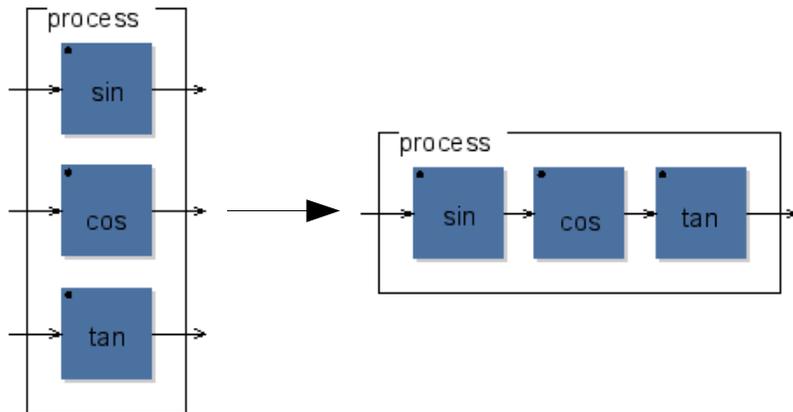
normal form

- Whitehead et al: *universal algebra*
- *Equational logic*: equality of normal forms
- O'Donnell et al: term rewriting as *programming language*
- Goguen, Mahr et al: *algebraic specification*
- Milner, Turner et al: *modern functional programming*

# Faust Term Rewriting Extension

```
serial((x,y)) = serial(x) : serial(y);  
serial(x)     = x;  
process      = serial((sin,cos,tan));
```

## BDA Term Rewriting



```
fold(1,f,x) = x(0);  
fold(n,f,x) = f(fold(n-1,f,x),x(n-1));  
fsum(n)     = fold(n,+);
```

```
f0 = 440; a(0) = 1; a(1) = 0.5; a(2) = 0.3;  
h(i) = a(i)*osc((i+1)*f0);  
v     = hslider("vol", 0.3, 0, 1, 0.01);  
process = v*fsum(3,h);
```

## Custom BDA Ops

# Faust Term Rewriting Extension

```
g(1,f) = f;
g(m,f) = (f, r(m-1)) : (_, g(m-1,f));

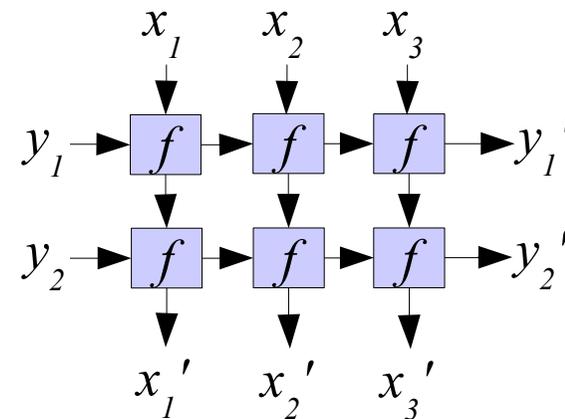
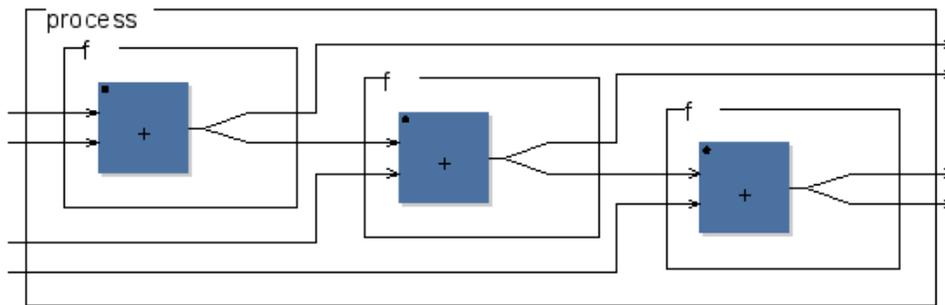
h(1,m,f) = g(m,f);
h(n,m,f) = (r(n+m) <: (!,r(n-1),s(m),
    (_,s(n-1),r(m) : g(m,f)))) :
    (h(n-1,m,f), _);

r(1) = _; r(n) = _,r(n-1); // route through
s(1) = !; s(n) = !,s(n-1); // skip

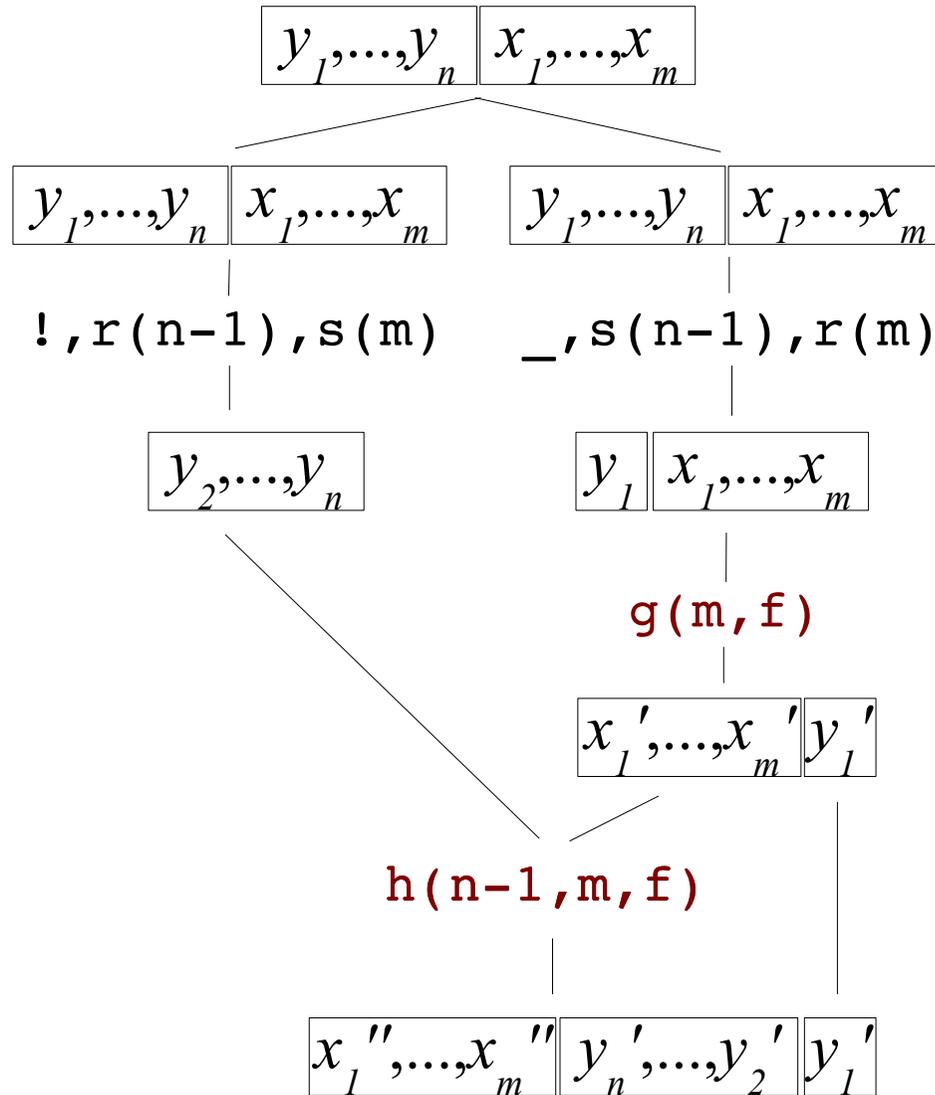
f = + <: _,_; // cell function
process = h(2,3,f);
```

**Systolic  
Array:**  
parallel  
processing  
in a 2D grid

$g(3,+)$ :



# Faust Term Rewriting Extension



**Systolic  
Array:**  
arranging  
the rows

$$\begin{aligned}
 h(n, m, f) &= (r(n+m) <: (!, r(n-1), s(m), \\
 &\quad (_, s(n-1), r(m) : g(m, f))) : \\
 &\quad (h(n-1, m, f), _));
 \end{aligned}$$

